

Differential Power Analysis of the HMAC Algorithm

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1. The HMAC Algorithm

- HMAC = keyed- $\underline{\mathbf{H}}$ ash $\underline{\mathbf{M}}$ essage $\underline{\mathbf{A}}$ uthentication $\underline{\mathbf{C}}$ ode
- Used in IPSec & TLS protocols for authentication & integrity
- Hash functions are not usually considered as targets for sidechannel attacks, as they are mostly used to process non-secret information. However, HMAC handles the secret key:

 $HMAC_{k}(x) = H((\mathbf{k} \oplus opad) || H((\mathbf{k} \oplus ipad) || x))$

- k = secret key ipad, opad = fixed public paddings x = message to be authenticated (variable)
- In the hash function *H*, **known variable** data mixes with **fixed unknown** data \rightarrow DPA attack is theoretically possible
- Attack would signatures for chosen messages to be forged
- Here, we present practical 1st-order attacks on HMAC-SHA-2 and HMAC-Whirlpool, and design countermeasures

2. Attacking HMAC-SHA-2

- SHA-2 family (Secure Hash Algorithm) is a well-known set of dedicated hash functions, standardised by NIST
- SHA-2 Compression Core:



- Goal of DPA attack on HMAC-SHA-2 is to recover fixed intermediate hash of (*k* ⊕ *ipad*), i.e.
 Fixed unknown data: initial states of registers A, B, C, ..., H Known variable data: message schedule *W*,
- A H can be recovered using seven DPA attacks [1] *However*, in SHA-256 the variables are 32-bit, and in SHA-512 the variables are 64-bit → Attack is difficult in practice
- Attack is simplified using **Partial Correlation** technique [2, 3]



given the correct prediction of D and the previous state of E.

Partial Correlation: Predict 32 bits

Predict *n* bits

Correlation = ρ Correlation = $\rho \sqrt{n/32}$

 \rightarrow Can make hypotheses on smaller sets of bits at a time

• Keep only those hypotheses with highest partial correlations (extend-and-prune approach)

 \rightarrow

• Build up partial correlations from $4 \rightarrow 32$ bits; much less computation than attempting immediate full 32-bit correlation.

3. Masking SHA-2



• Masked circuits designed for *Ch* and *Maj* functions [1]

• Boolean-to-Arithmetic and Arithmetic-to-Boolean conversions optimised for FPGA using dedicated **carry chain**

4. HMAC-Whirlpool: Attack & Masking

- Whirlpool hash function recommended in *NESSIE* portfolio. Block cipher-based, similar to AES
- Similar attack goal to HMAC-SHA-2 case: recover secret intermediate chaining state
- Attack is less complex due to Whirlpool S-box: Processes each byte of the 512-bit state independently
 → Can focus on 8-bit intermediate variables in the DPA attack
- Masked circuit [4]:



- Masked S-box S' = 5 masked 4-bit pre-computed look-up tables. Re-use of round function transformations to compute mask correction → negligible area impact
- **Future work**: Consider template attacks & higher order attacks Future hash functions? (NIST Hash Function Competition)

References

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Cryptographic Hardware and Embedded Systems - CHES 2007, 9th International Workshop, Vienna, Austria, September 2007